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A Competent Approach for Protected Communication with Double EHDES (ACAPCD-E)

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Abstract: In this paper, we established a secure protocol of communication where EHDES provides a high level security on e-mail Transmission.

Keywords: Message, EHDES, Double EHDES, Steganography, Covert Mailing System.

I. INTRODUCTION

Today's e-mail is more important way of communication between anonymous user and within an organization. Steganography has a relatively short history; even today ordinary dictionaries do not contain the word "steganography". Books on steganography are still very few [1], [2]. The very large data hiding capacity is the main feature of steganography [3], [4]. Some information systems use steganography for hiding the data. Some key is used in these systems when it embeds/extracts secret data. One natural application is a secret mailing system [5], [6] that uses a symmetric key. Another application pays attention to the nature of steganography whereby the external data (e.g., visible image data) and the internal data (any hidden information) cannot be separated by any means. We will term this nature as an "inseparability" of the two forms of data.

In this proposed scheme, we will show an example of a mixed scheme of steganography and cryptography is, A Competent Approach for Protected Communication with Double EHDES, which isan anonymous and covert e-mailing system with complete security.

This paper is described in different steps. In Step 2,we describes the scheme of double enhanced data encryption standard (D-EHDES).In Step 3, we will show a competent approach for protected communicationusing symmetric key. How we can make it a safe system described in Section 4. Finally, section 5 is conclusion.

II. PRELIMINERIES

The amount of transfer messaging has increased rapidly on the Internet. Cryptography is a branch of applied mathematics that aims to add security in the ciphers of any kind of messages. Cryptography algorithms use encryption keys, which are the elements that turn a general encryption algorithm into a specific method of encryption. The data integrity aims to verify the validity of data contained in a given document. [7]

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A. Double EHDES

Double EHDES is an arrangement or cascading of EHDES and its working just like a EHDES but two times. In Enhanced Data Encryption Standard (EHDES) [8, 9], we breaks block of message and follow these three phases: 1. Key Generation. 2. Encryption.

3. Decryption.

Data or Message (M)	56 Bit Symmetr ic Key 'k'		56 Bit Symmet ric Key 'k'	
	EHDES		EHDES	Data or Message (M)
M ₁ , M ₂ , M ₃ Mn		Cipher text or Encrypted Message (C)		M ₁ , M ₂ , M ₃ Mn
	Encryption Process		Decryption Process	

[a] Key Generation

In this phase, EHDES generates the n different keys $(K_{new1}, K_{new2}, K_{new3}, \dots, K_{new n})$ to apply a function F on Initial key k and a random number (N_{RNG}) , for every block of message (M_1, M_2, M_3, M_n) .

[b] Encryption on Input Data.

Message breaks in 64 Bit n blocks of plain text. $M = \{M_1, M_2, M_3, \dots, M_n\}$ Now, we encrypt our message $\{M_1, M_2, M_3, \dots, M_n\}$ blocks by each new generated key $K_{new1}, K_{new2}, K_{new3}, \dots, K_{new n}$.

[c] Decryption on Input Cipher:

Decryption is the reverse process of encryption. For decryption, we also used the same key which is used in encryption. On the receiver side, the user also generate the same new key $K_{\text{new i}}$ for each block of cipher and generate plain text through decryption process of data encryption standard.

III. A MODEL OF PROTECTED COMMUNICATION ACAPCD-E

A Competent Approach for Protected Communication with Double EHDES (ACAPCD-E) is a steganography application program with cryptography. In the following description, $M_{E_{ACAPCD-E}1}$, denotes a member of ACAPCD-E 1, and $M_{E_{ACAPCD-E}2}$, denotes a member of ACAPCD-E 2.

An ACAPCD-E consists of the three following components.

- [a] Envelope Producer (EP).
- [b] Message Inserter (MI).
- [c] Envelope Opener (EO).

We denote $M_{E_{ACAPCD-E}I}$'s ACAPCD-E as $ACAPCD - E_I$ (i.e., customized ACAPCD-E by $M_{E_{ACAPCD-E}I}$. So, it is described as $M_{E_{ACAPCD-EI}} = (EP_{E_{ACAPCD-EI}}, MI_{E_{ACAPCD-EI}}, EO_{E_{ACAPCD-EI}})$ $EP_{E_{ACAPCD-E}I}$ is a component that produces $M_{E_{ACAPCD-E}I}$'s envelope $(E_{E_{ACAPCD-EI}})$ and $af = \sum_{i=1}^{n} i$. $E_{E_{ACAPCD-EI}}$ is the envelope (actually, an image file) which is used by all other members in the organization when they send a secret message to $M_{E_{ACAPCD-E}I}$. $(EO_{E_{ACAPCD-E}I})$ is produced from an original image (EO). $M_{E_{ACAPCD-E^{I}}}$ can select it according to his preference. $(E_{E_{ACAPCD-EI}})$ has both the name and e-mail address of $M_{E_{ACAPCD-E}I}$ on the envelope surface (actually, the name and address are "printed" on image $(E_{E_{ACAPCD-E}I})$. It will be placed with function f at an open site in the organization so that anyone can get it freely and use it any time. Or someone may ask $M_{E_{ACAPCD-E}I}$ to send it directly to him/her. $(MI_{E_{ACAPCD-EI}})$ is the component to insert (i.e., embed according to the stegnographic scheme) $M_{E_{ACAPCD-EI}}$'s message into another member's (e.g., $M_{E_{ACAPCD-E2}}$)'s envelope $(E_{E_{ACAPCD-E^2}})$ when $M_{E_{ACAPCD-E^1}}$ is sending a secret message $(Mess._{E_{ACAPCD-E1}})$ to $(M_{E_{ACAPCD-E2}})$. One important function of $M_{E_{ACAPCD-E}1}$ is that it detects a key $(Key_{E_{ACAPCD-E}1})$ that has been hidden in the envelope($E_{E_{ACAPCD-E^2}}$), and uses it when inserting a message ($Mess._{E_{ACAPCD-E^1}}$) in ($E_{E_{ACAPCD-E^2}}$). $(EO_{E_{ACAPCD-E}1})$ is a component that opens (extracts) "message $(E_{E_{ACAPCD-E}1})$'s inserted" envelope $(E_{E_{ACAPCD-E^1}}(Mess_{E_{ACAPCD-E^2}}))$ which $M_{E_{ACAPCD-E^1}}$ received from someone as an e-mail attachment. The sender $(M_{E_{ACAPCD-E^2}})$ of the secret message $(Mess._{E_{ACAPCD-E^2}})$ is not known until $M_{E_{ACAPCD-E1}}$ opens the envelope by using $(EO_{E_{ACAPCD-E1}})$.

IV. CUSTOMIZATION OF AN ACAPCD-E

Customization of an ACAPCD-E for member $(M_{E_{ACAPCD-E1}})$ takes place in the following way. $(M_{E_{ACAPCD-E1}})$, © 2010, IJARCS All Rights Reserved

first decides a key $(Key_{E_{ACAPCD-E^1}})$ with $f = \sum_{i=1}^n i$ where i is a positive integer, when he/she installs the ACAPCD-E onto his computer. Let us suppose $E_{ACAPCD-E}$ 2 try to communicate at any time t, then he/she picks up a number randomly form i. Now, ACAPCD-E generates $f_t = \sum_{i=1}^{n-1} i$. Let $R = f - f_t$, ACAPCD-E generate a key $(Key_{E_{ACAPCD-E1}})$ with the help of R using Double EHDES key generation process. Then he types in his name e-mail $(Name_{E_{ACAPCD-E}1})$ and address $(Emailadr_{E_{ACAPCD-E}1}).(Key_{E_{ACAPCD-E}1})$ is secretly hidden (according to a steganographic procedure in his envelope $(E_{E_{ACAPCD-E1}})$. This $(Key_{E_{ACAPCD-E1}})$ is eventually transferred to a message sender's $(MI_{E_{ACAPCD-E^2}})$ in an invisible way. $(Name_{E_{ACAPCD-E}1})$ and $(Emailadr_{E_{ACAPCD-E}1})$ are printed out on the envelope surface when $(M_{E_{ACAPCD-E}1})$ produces $(E_{E_{ACAPCD-E1}})$ by using $(EP_{E_{ACAPCD-E1}})$. $(Key_{E_{ACAPCD-E1}})$ is also set to $(EO_{E_{ACAPCD-E}1})$, when communicators wish to start the communication. $(Name_{E_{ACAPCD-E^1}})$ and $(Emailadr_{E_{ACAPCD-E^1}})$ are inserted (actually, embedded) automatically by also $(MI_{E_{ACAPCD-E}1})$ any time $(M_{E_{ACAPCD-E}1})$ inserts his message $(Mess._{E_{ACAPCD-E1}})$ in another member's envelope $(E_{E_{ACAPCD-E2}})$. The embedded $(Name_{E_{ACAPCD-E}1})$ and $(Emailadr_{E_{ACAPCD-E}1})$ are extracted by a message receiver $(M_{E_{ACAPCD-E^2}})$ by $(EO_{E_{ACAPCD-E^2}})$.

A. How it works:

When some member $(M_{E_{ACAPCD-E^2}})$ wants to send a secret message (Mess._{EACAPCD-E2}) to another member ($M_{E_{ACAPCD-E1}}$), whether they are acquainted or not, $(M_{E_{ACAPCD-E}2})$ gets (e.g., downloads) the $(M_{E_{ACAPCD-E1}})$'s envelope $(E_{E_{ACAPCD-E1}})$, and uses it to insert his message (Mess._{EACAPCD-E2}) by using (MI_{EACAPCD-E2}). When $(M_{E_{ACAPCD-E^2}})$ tries to insert a message, $(M_{E_{ACAPCD-E^1}})$'s key $(Key_{E_{ACAPCD-E1}})$ is transferred to $(MI_{E_{ACAPCD-E2}})$ automatically in an invisible manner, and is actually used. $(M_{\text{SES}_{\text{EHDES}}I})\,\text{Can}$ send $(E_{E_{ACAPCD-E1}}(M_{E_{ACAPCD-E2}}))$ directly, or ask someone else to send, it to $(M_{E_{ACAPCD-E}1})$ as an e-mail attachment with using encryption process of Double EHDES. $(M_{E_{ACAPCD-E^2}})$ Can be anonymous information sender's is because no seen on $(E_{E_{ACAPCD-E1}}(M_{E_{ACAPCD-E2}}))$. (Mess. $E_{ACAPCD-E2}$) is hidden, and only $(M_{E_{ACAPCD-E}1})$ can see it by opening the envelope. It is not a problem for $(M_{E_{ACAPCD-E^2}})$ and $(M_{E_{ACAPCD-E^1}})$ to be acquainted or not because $(M_{E_{ACAPCD-E^2}})$ can get anyone's envelope from an open site.

V.CONCLUSION

This paper provides the effective solution of complete and high level security of an e-messaging system. Steganography and the key generation process of Double-EHDES is the strength of this model. ACAPCD-E is a very easy-to-use system because users are not bothered by any key handling, as the key is always operated automatically. As ACAPCD-E doesn't need any authorization bureau.

Table I

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STEP	ENTITY	PROCESS	
1.	ACAPCD-E 1	A. Generate an envelope $E_{ACAPCD-E} 1$. B. Upload an envelope $E_{ACAPCD-E} 1$ and X (set of positive integer) with $f = \sum_{i=1}^{n} i$. C. Name Name _{ACAPCD-E} 1 and Email address Emailadr _{ACAPCD-E} 1 print on envelope $E_{ACAPCD-E} 1$ is surface and key Key _{ACAPCD-E} 1 is hiding using stegnohraphic function in the envelope $E_{ACAPCD-E} 1$.	
2.	ACAPCD-E 1	Choose any number i randomly fromX.	
3.	Downloadable Site Function	A. Calculate $f_t = \sum_{i=1}^{n-1} i$ and $R = f - f_t$. B. Send R to ACAPCD - E1.	
4.	Downloadable Site Function	 A. Moderate key Key_{ACAPCD-E}1 using EHDES key process. B. Update keyKey_{ACAPCD-E}1. 	
5.	ACAPCD-E 2	Download envelope $E_{ACAPCD-E}$ 1.	
6.	ACAPCD-E 2	 A. Picked out key Key_{ACAPCD-E}1 from envelopeE_{ACAPCD-E}1. B. Insert message Mess_{-ACAPCD-E}2 in the envelopeE_{ACAPCD-E}1. C. Encrypt the envelope contains message(E_{ACAPCD-E}1 (Mess_{-ACAPCD-E}2))using Double EHDES with key Key_{ACAPCD-E}1. 	
7.	ACAPCD-E 2	Send envelope contains message($E_{ACAPCD-E}1$ (Mess. _{ACAPCD-E} 2)) toACAPCD - E1.	
8.	ACAPCD-E 1	 A. Receive envelope contains message(E_{ACAPCD-E}1 (Mess._{ACAPCD-E}2)). B. Decrypt the envelope contains message(E_{ACAPCD-E}1 (Mess._{ACAPCD-E}2))using Double EHDES with key Key_{ACAPCD-E}1. C. Separate messageMess._{ACAPCD-E}2 from the envelopeE_{ACAPCD-E}1. D. Read messageMess	

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