

THE $(n-1)/2$ -REGULAR GRAPH ON n VERTICES

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Abstract. Let G be an undirected and simple graph on n vertices and degree of each vertex is equal $(n-1)/2$. We present some properties of G and confirm that G is a Hamiltonian graph.

Keywords. Regular graph, Hamiltonian graph, Petersen graph, Closure graph, Diameter of Graph

1. INTRODUCTION

Let $G = (V, E)$ be an undirected and simple graph on n vertices, where V be the vertex set and E be edge set of G . We use $|V|$ and $|E|$ to denote the number of vertices and the number edges of G , respectively. In G , the degree of vertex v is denoted by $\deg(v)$. The edge of two vertices u and v is denoted by (u, v) or uv . A graph is called *regular graph of degree k* (or *k -Regular graph*) if its vertices has degree k . We use $\delta(G)$ to denote the minimum degree of the vertices of G . The graph on n vertices with all vertices having degree $n-1$ is called the *complete graph* and denote by K_n .

A set of vertices in graph G is called *independent* if no two vertices in this set are non-adjacent. *Maximum independent set* is an independent set of largest possible size for a given graph. Denote by $\alpha(G)$ the size of a maximum independent set of G . A set $C \subseteq V$ is called *clique* if every two distinct vertices in C are adjacent in G .

The graph $H = (W, F)$ is called a *subgraph* of G if $W \subseteq V$ and $F \subseteq E$. Let v is a vertex of G , we use $G-v$ to denote the subgraph which obtained by deleting v from G . Livewise, if B is a set of vertices of G , graph $G-B$ is a subgraph of G whose obtained by deleting B from G .

We use $\omega(G)$ to denote the number of components of G . In G , a vertex v is called *cut vertex* if $\omega(G) < \omega(G-v)$. Denote by $G+uv$ the graph which obtained from G when previously non-adjacent vertices u

and v are joined by a new edge uv . A set of vertices in a connected graph is called *disconnecting* if the graph becomes disconnected when this set is removed. Denote by $\kappa(G)$ the smallest size of a disconnecting set in G .

Graph G is called *1-tough* if $\omega(G-B) \leq |B|$ for every non-empty subset B of V .

The distance between two vertices in G is the number of edges in a shortest path connecting them. The *diameter* of G is the greatest distance between any pair of vertices and denote by $d(G)$.

A simple path in connected graph G that passes through every vertex exactly once is called *Hamiltonian path*. A simple cycle in a connected graph G that passes through every vertex exactly once is called *Hamiltonian cycle*. Any connected graph that contains a Hamiltonian cycle is called *Hamiltonian Graph*.

Recognizing Hamiltonian graph is hard problem. Now there are many theorems providing sufficient conditions for a graph to be Hamiltonian. Dirac [4] proved that if the minimum degree of the vertices of G is at least $n/2$ then G is Hamiltonian graph. Denote by $\sigma_2(G)$ - the degree sum of any two non-adjacent vertices in G . Ore [4] asserts results more generally, if $\sigma_2(G) \geq n$ then G is Hamiltonian graph. In [4], H. A. Jung proved that, if G is 1-tough and $\sigma_2(G) \geq n-4$, $n \geq 11$ then G is Hamiltonian graph.

In [1] and [2], we proved that, if $\sigma_2(G) = n-1$, there are three cases, if n is an even number then G is Non-Hamiltonian graph, if n is an odd number and

$2 < \alpha(G) < (n+1)/2$ then G is Hamiltonian graph, otherwise, G is Non-Hamiltonian graph.

In [5], Paul Erdos proved that, if $(n-2)$ -Regular G graph with $|V(G)|=2n$ or $|V(G)|=2n-1$ and $\kappa(G)=2$, then, G is Hamiltonian if only if G is not the Petersen graph. Figure 1 is Petersen graph.

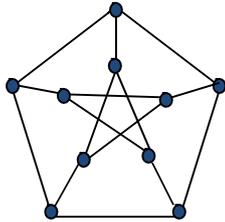


Figure 1. Petersen graph.

Bondy, Chvátal and Murty [3] used the definition on closure graph to define the necessary and sufficient condition for Hamiltonian graph. Following some sufficient conditions for Hamiltonian and non-Hamiltonian graph.

Theorem 1 (Bondy and Chvátal [3]). *Let G be a graph on n vertices and let u and v be nonadjacent vertices of G with degree sum at least n . Then, G is Hamiltonian graph if and only if $G+uv$ is Hamiltonian graph.*

Theorem 2 (Chvátal [3]). *If G is not 1-tough graph then G is not Hamiltonian graph.*

Denote by $Cl(G)$ the closure of G which derived from G by recursively joining pairs of nonadjacent vertices having degree sum at least n . Figure 2 illustrates graph G and its closure graph $Cl(G)$.

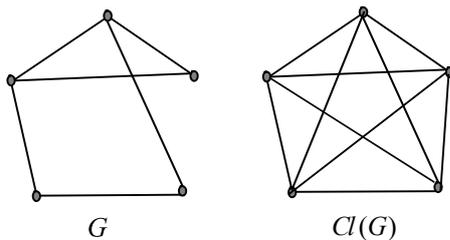


Figure 2

Theorem 3 (The Closure Lemma). *G is Hamiltonian if and only if $Cl(G)$ is Hamiltonian.*

Following result is special case of Theorem 3.

Corollary 1 (Bondy and Murty [3]). *If $Cl(G)$ is complete graph K_n then G is Hamiltonian.*

Theorem 4 (Nash-Williams, Bondy [5]). *If $\alpha(G) \leq \delta(G)$, $\kappa(G) \geq 2$ and $\delta(G) \geq (n+2)/3$ then G is Hamiltonian.*

2. RESULT

Let G be an k -regular graph on n vertices, where $k = (n-1)/2$. Then, n must be an odd number and $\text{mod}(n-1,4) = 0$ (if not, $(n-1)/2 = k$ be an odd number, i.e., graph G has number of vertices of odd degree is an odd number, this is absurd).

We use $G(n,k)$ to denote the set of k -regular graphs on n vertices, where $k = (n-1)/2$ and $\text{mod}(n-1,4) = 0$ (so, $n \geq 5$ and k be an even number). Figure 3 illustrates graphs in $G(5,2)$ and $G(9,4)$.

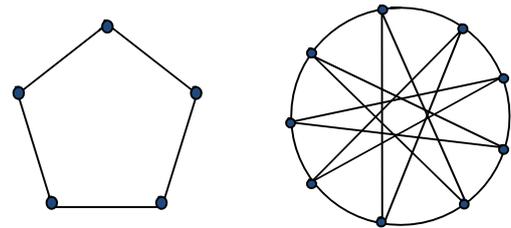


Figure 3. Graphs in $G(5,2)$ and $G(9,4)$.

Proposition 1. *For every $G \in G(n,k)$, G is connected graph.*

Proof. Suppose otherwise, G is disconnected graph. Let G^1 is a connected component of G and $|V(G^1)|=n_1$. Denote by G^2 the remaining of G and $|V(G^2)|=n_2$. We have $n_1+n_2=n$. Choose an any vertex u in G^1 and an any vertex v in G^2 . Then, $(n-1)/2 = \text{deg}(u) \leq n_1-1$, $(n-1)/2 = \text{deg}(v) \leq n_2-1$. So, $n-1 = \text{deg}(u) + \text{deg}(v) \leq n_1-1+n_2-1 = n-2$, a contradiction. Therefore, G is connected graph.

Proposition 2. *For every $G \in G(n,k)$, G contains a Hamiltonian path.*

Proof. Let u and v be any two non-adjacent vertices in G , we add an edge uv to G . Then, $\text{deg}(u) = \text{deg}(v) = 1+(n-1)/2$. Let w is an any vertex such that w is non-adjacent to u or v of G , we have $\text{deg}(w) + \text{deg}(u) = (n-1)/2 + 1 + (n-1)/2 = n$ or

$\deg(w) + \deg(v) = (n-1)/2 + 1 + (n-1)/2 = n$. In other words, we add to the $G+uv$ graph the edges connecting two non-adjacent vertices whose degree sum is not less than n . Thus, $Cl(G+uv)$ is complete graph K_n , and by Corollary 1, $G+uv$ is Hamiltonian graph. This proves that, G contains a Hamiltonian path.

Note that, for $n=5$, $G(5,2)$ has only one graph as shown in Figure 3.

Suppose that, $G \in G(n,k)$, u and v are two non-adjacent vertices in G . Denote by N_v the set of vertices that are non-adjacent to v , N_u the set of vertices that are non-adjacent to u in G . Thus, $Z = V \setminus N_u \cup N_v$ is a set of vertices which are both adjacent to v and u , $A = N_u \cap N_v$ is a set of vertices which are non-adjacent to v and u .

Proposition 3. For every $G \in G(n,k)$, $|Z| = |A| + 1$.

Proof. By all vertices of the G have degrees $(n-1)/2$, $|N(u)| = n-1 - \deg(u) = n-1 - (n-1)/2 = (n-1)/2$.

Similarly, $|N(v)| = (n-1)/2$. We have, $|Z| = |V \setminus [N_u \cup N_v]| = n - [|N_u| + |N_v| - |N_u \cap N_v|] = n - [(n-1)/2 + (n-1)/2 - |A|] = |A| + 1$. Thus, $|Z| = |A| + 1$.

Proposition 4. For every $G \in G(n,k)$, $d(G) = 2$.

Proof. Let u and v be two non-adjacent vertices in G . By Proposition 3, $|Z| = |A| + 1$, so $|Z| \geq 1$, or $Z \neq \emptyset$. This proves that, with two non-adjacent vertices u and v in G , there exists at least one vertex $z \in Z$ such that z is adjacent to both vertices u and v . In other words, $\forall (u,v) \notin E(G)$, $d(u,v) = 2$. Thus, $d(G) = 2$.

Proposition 5. Let $n \geq 9$, for every $G \in G(n,k)$, $3 \leq \alpha(G) \leq (n-1)/2$.

Proof. a) First, we will prove that $3 \leq \alpha(G)$.

Assume that $\alpha(G) = 2$. Let u and v be two any non-adjacent vertices in G .

Consider 1. By $\alpha(G) = 2$, so $A = \emptyset$, and by Proposition 3, $|Z| = 1$. Let $Z = \{z\}$, and so z is the only vertex that is adjacent to both vertices u and v in G . Let N_{uz} be the set of vertices of N_u that are non-adjacent to z

, N_{vz} be the set of vertices of N_v that are non-adjacent to z . Figure 4 illustrates a graph in $G(9,4)$ to prove Proposition 5.

Figure 4.

Obviously, $|N_{vz}| + |N_{uz}| = (n-1)/2$. Moreover, by $\alpha(G) = 2$, each pair of vertices in N_{uz} must be adjacent, and each vertex in N_{uz} must be adjacent to every vertex in N_{vz} . Similarly, each pair of vertices in N_{vz} must be adjacent. In other words, the vertices in N_{uz} form a clique $K_{|N_{uz}|-1}$ and the vertices in N_{vz} form a clique $K_{|N_{vz}|-1}$ in G .

Consider 2. Suppose that w is any vertex in N_{vz} . Then, there exists at least one vertex $r \in N_v \setminus N_{vz}$ such that w is adjacent to r (if not, graph G will have three vertices w, r, v , where each pair is non-adjacent, is contradictory to hypothesis $\alpha(G) = 2$).

From *Consider 1* and *Consider 2*, we have, vertex w must be adjacent to u, r and all vertices in N_{uz} and N_{vz} . I.e., $\deg(w) \geq 1 + 1 + |N_{vz}| - 1 + |N_{uz}| = 1 + (n-1)/2$. This is contrary to the assumption of the k -regular graph G , $k = (n-1)/2$. So, $\alpha(G) \geq 3$.

b) Next, we will prove that $\alpha(G) \leq (n-1)/2$.

Assume that $\alpha(G) = (n+1)/2$, and let $S = \{s_1, s_2, \dots, s_{(n+1)/2}\}$ is a maximum independent set of G . Set $M = V \setminus S$. We have, $|M| = n - |S| = n - (n+1)/2 = (n-1)/2$. For every $i \in \{1, 2, \dots, (n+1)/2\}$, $\deg(s_i) = (n-1)/2$. So s_i is adjacent to $(n-1)/2$ vertices in M . I.e., each vertex in M must be adjacent to every vertex in $S = \{s_1, s_2, \dots, s_{(n+1)/2}\}$. This proves that, each vertex in M has degree no less than $(n+1)/2$, this is contrary to the

assumption of the k -regular graph G . Therefore, $\alpha(G) \leq (n-1)/2$.

Note that, Proposition 5 is also true for $n=5$, in $G(5,2)$ has the only graph G for $\alpha(G) = (5-1)/2 = 2$ (see Figure 3). Figure 5 illustrates graphs in $G(9,4)$ for $\alpha(G) = 3$ and $\alpha(G) = 4$.

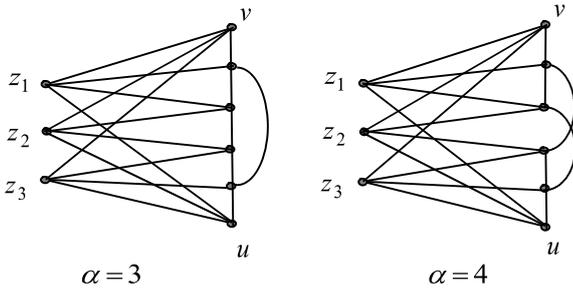


Figure 5.

Theorem 5. Let $n \geq 9$, for every $G \in G(n,k)$, G is Hamiltonian graph.

Proof. We will show that graph G satisfies the condition of Theorem 4, and therefore, Theorem 5 is proved.

Indeed, by $\delta(G) = k = (n-1)/2$ (the hypothesis of G) and $3 \leq \alpha(G) \leq (n-1)/2$ (Proposition 3), so $\alpha(G) \leq \delta(G)$. (1)

By $n \geq 9$, we have $(n-1)/2 \geq (n+2)/3$, i.e., $\delta(G) \geq (n+2)/3$. (2)

Next, we show that $\kappa \geq 2$. Suppose otherwise, $\kappa = 1$ and w is an any cut vertex of G . Then, graph $G-w$ is disconned graph, and in $G-w$ there exist two disjoint sets X and Y such that $V = \{w\} \cup X \cup Y$, $X \cap Y = \emptyset$. By, each vertex in G has degree $\delta = (n-1)/2$, so $|X| = |Y| = (n-1)/2$, all vertices of X (similarly Y) whose each pairwise are adjacent, and all vertices of $X \cup Y$ are adjacent to w . So, $\deg(w) = |X| + |Y| = (n-1)/2 + (n-1)/2 = n-1$, a contradiction with the hypothesis of G . Thus, $\kappa \geq 2$. (3)

From (1), (2), (3) shown that graph G satisfies the condition of Theorem 4.

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