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Incorporating 3G Networks in Wireless Environments

B.V.S.S.R.S.Sastry* Aurora's Engineering College Bhongir, Andhra Pradesh, India sastry_38@yahoo.com S.Santhi Priya Asst. Professor (IT), Aurora's Engineering College Bhongir, Andhra Pradesh, India santhipriya.sunkara@gmail.com

K. Akshitha Aurora's Engineering College Bhongir, Andhra Pradesh, India koluguri.87@gmail.com

Abstract: The third generation (3G) cellular networks provide ubiquitous connectivity but low data rates, whereas wireless local area networks (WLANs) can offer much higher data rates but only cover smaller geographic areas. Their complementary characteristics make the integration of the two networks a promising trend for next-generation wireless networks. With combined strengths, the integrated networks will provide both wide-area coverage and high-rate data services in hotspots. There are many aspects involved in their interoperability, such as mobility, security, and quality of service (QoS) provisioning. In this article, we present a survey of most recent interoperability mechanisms proposed in the literature, and outline some important open issues to achieve seamless integration.

Keywords: 3G/WLAN interoperability; tight/loose coupling; mobility management; QoS provisioning.

I. INTRODUCTION

In the past decade, there has been successful deployment and fast evolution of various wireless networks. Different technologies are adopted for different application environments. It is well recognized that the next-generation wireless networks will integrate heterogeneous technologies to achieve enhanced performance. The attractive and complementary characteristics presented by cellular networks and wireless local area networks (WLANs) make them promising candidates. Originally aiming at providing high-quality circuit- switched voice service to mobile users within wide areas, cellular networks have been well deployed around the world and have evolved to the third generation (3G) [1]. Two major standards for 3G mobile/wireless cellular networks are the universal mobile telecommunication system (UMTS) and cdma2000, which are specified by the 3rd generation partnership projects, i.e., 3GPP and 3GPP2, respectively.

Both systems are based on code-division multiple access (CDMA) and augmented with packet-switched data services, such as multimedia message service (MMS), and wireless application protocol (WAP) service. In cdma2000, for example, the nominal 1.25 MHz bandwidth can achieve a data rate up to around 2 Mbps for indoor office environments. However, this is still not enough to satisfy the ever-increasing demands for bandwidth-intensive data applications.

On the other hand, usually operating at unlicensed frequency bands, WLANs provide data services with lower cost. Moreover, the large bandwidth available for WLANs makes it possible to achieve higher data rates. For example, in IEEE 802.11b (one of the most popular WLAN standards), a WLAN can have a bandwidth more than 20 MHz. IEEE 802.11b operates at the license-exempt industrial, scientific and medical (ISM) frequency band from 2.4 GHz to 2.483 GHz. It extends the physical layer based on direct sequence spread spectrum (DSSS) specified in the original 802.11 standard and supports a higher data rate up to 11 Mbps. The subsequent revisions such as 802.11a and 802.11g adopt orthogonal frequency-division multiplexing (OFDM) and offer a maximum rate of 54 Mbps at the unlicensed 5 GHz and 2.4 GHz bands, respectively. However, designed as a wireless

extension to the wired Ethernet, a WLAN can only cover a small geographic area. For instance, an 802.11b access point (AP) can communicate with a mobile station (MS) within up to 60 m at 11 Mbps and up to 100 m at 2 Mbps with omnidirectional antennas.

Consequently, with lower cost and much higher data rates, WLANs can effectively supplement the 3G networks in hotspot areas, where bandwidth-demanding applications are concentrated. As a result, by effectively combining 3G cellular networks and WLANs into an integrated wireless data access environment, mobile users can be provided with both ubiquitous connectivity and high-rate data services in hotspots.

In the following sections, we first discuss the important challenging issues involved in this integration problem. Then we briefly review some typical interoperability solutions proposed in the literature.

II. CHALLENGES FOR 3G / WLAN INTEROPERABILITY

The heterogeneous technologies employed in 3G cellular networks and WLANs bring many challenges to the interoperability. Based on different radio access techniques, cellular networks and WLANs present distinct characteristics in terms of mobility management, security support, and quality of service (QoS) provisioning. In order to achieve seamless integration, these issues should be carefully addressed while developing the interoperability schemes.

After three-generation evolution, relatively mature and complete technologies have been established in cellular networks to address issues such as mobility, security, QoS, etc. With widely-deployed infrastructure from radio access networks to core networks, ubiquitous connectivity is provided to mobile uses over wide areas. Different mobility levels are supported from fast vehicles moving on highways to stationary users in an indoor environment. In contrast, the WLAN specifications only focus on the physical layer and medium access control (MAC) layer. As for the upper layers, it assumes to adopt the same protocols as those in wired networks, e.g., the Internet protocol (IP) suite, with some adaptation for wireless links to avoid performance degradation. A de facto WLAN system is given in [1], in which the layer-2 distribution system connects multiple APs, while access routers in

turn connect the layer-2 distribution system to an ${\rm IP}$ backbone network.

In some proposed interoperability schemes (to be discussed), the access routers offer rich functionalities more than the basic function of IP routing, e.g., transferring authentication and charging information between the 3G networks and WLANs. Through border gateways in the IP backbone network, WLAN terminals are provided IP connectivity to external IP networks such as the public Internet or a corporation intranet. Instead of providing continuous coverage over wide areas, WLANs are usually disjointly deployed in public or private hotspots such as cafes, airports, and offices. Users in these areas normally have a very low mobility level, as most of these areas are located in indoor environments. Also, cellular coverage is available in these areas. As a result, a non-uniform overlay topology structure has to be considered for 3G/WLAN integration.

A. Seamless Roaming across 3G Cellular Networks and WLANs

Taking into account the distinct mobility management mechanisms employed in cellular networks and WLANs, it is a rather challenging task to support seamless roaming across the two networks. Either the cellular networks or WLANs should have inherent mechanisms for location and handoff management to support the layer-2 or link layer mobility. In 3G networks (e.g., UMTS or cdma2000), with the aid of core networks, tunnelling protocols are used to support roaming within the public land mobile network (PLMN) or across 3G PLMNs of different operators with roaming agreements. For example, in UMTS, the general packet radio service (GPRS) mobility management (GMM) is used. However, link layer mobility is not enough to provide network-layer transparency to upper-layer applications. To avoid disruption of upper-layer sessions due to IP address changes when user mobility results in changes of network attachment points, the network-layer IP mobility is needed. In cdma2000, Mobile IP (MIP) [2] provides IP mobility within the same packet data serving node (PDSN) and between different PDSNs. However, in current specifications of UMTS, IP mobility is only supported within a UMTS network with the same gateway GPRS support node (GGSN). It is being considered to introduce Mobile IP for inter-UMTS or inter-technology IP mobility through a three-state evolution in [3]. An overview of the mobility management in UMTS and cdma2000 can be found in [4]. The mobility management in WLANs is much simpler since they are only oriented to local areas. In IEEE 802.11 WLANs, the distribution system (e.g., an 802.3-type Ethernet) connects multiple basis service sets (BSSs) into an extended service set (ESS). Each BSS is under control of an AP in the infrastructure mode. In this case, mobility across the BSSs within an ESS is handled by the APs involved. The inter-access point protocol (IAPP) specified in 802.11f further facilitates the user roaming between APs of different vendors. When IP connectivity is provided in the WLAN system, IP micro-mobility protocols can be introduced to further support IP layer mobility. Keeping in mind the difference of cellular networks and WLANs in mobility management, to achieve seamless roaming across the two networks, either a unified mobility management mechanism is followed, or both networks adopt their individual mobility management mechanisms and maintain proper interoperation.

B. Enhanced Security Level

Network security covers diverse issues such as user authentication, data confidentiality and integrity, and key management. The security mechanisms in 3G networks are built upon those used in the second-generation cellular networks, e.g., the global system for mobile communications (GSM) for the UMTS system. In particular, user authentication in UMTS adopts the authentication and key agreement (AKA) procedure, which relies on the universal subscriber identity module (USIM) running in a smart card at the user terminal. In addition to authenticating the subscriber's identity, cipher and integrity session keys are also generated from the long-term pre-shared secret key stored in the

USIM module and home location register/authentication center (HLR/AuC). A good introduction to the access security in 3G networks is given in [5, 6]. In the original 802.11 standard, rather weak security is provided due to lack of key management and flaws of the wired equivalent privacy (WEP) protocol. As a consequence, it fails the claimed objectives with respect to user identity privacy and data confidentiality. The 802.11i standard aims at improving the security level of WLANs. Port based access control standard IEEE 802.1X is introduced to enhance authentication and key management using the extensible authentication protocol (EAP) [7]. Also, as an interim solution, temporal key integrity protocol (TKIP) is specified to _x the vulnerabilities of WEP, while the long- term solution is based on advanced encryption standard (AES) in place of the stream cipher RC4 used in WEP. More details on 802.11 security are given in [8]. As we know, a system is only as strong as its weakest link. For 3G/WLAN interworking, if comparable security cannot be provided by both networks, adversaries can break into the system through the weakest component in the security chain and in turn defeat all security goals of the entire system. In other words, the two networks must be integrated in such a manner to achieve an enhanced (instead of impaired) security level. Also, appropriate independence between them should be maintained to minimize the security exposure and Domino effect when one of them is broken.

C. Consistent End-to-end QoS Guarantee

The shared nature of radio link necessitates proper medium access control to coordinate multiple connections to access the shared wireless channel. Both UMTS and cdma2000 use CDMA for multiple access. Moreover, by introducing packet-switched mode a higher resource utilization can be achieved for bursty data traffic by statistical multiplexing. For the multiple access uplink (from MS to base station), a two-phase request-grant access procedure is used in 3G networks. First MSs send transmission requests to the base station (BS) through a contention channel. The BS acknowledges those successful requests and reserves resources for data transmission to follow. Then the MSs are notified the resource assignments. This type of centralized control and reservation-based resource allocation, together with proper admission control to limit the traffic load, enables fine QoS provisioning in 3G cellular networks. On the other hand, the original WLAN specifications only support best-effort data service with contention-based random access protocol, e.g., the distributed coordination function (DCF) of 802.11 WLAN based on carrier sense multiple access/collision avoidance (CSMA/CA). In this mode, the AP competes for access with MSs instead of scheduling the resource assignments as a BS in 3G networks. The other centralized control mode, point coordination function (PCF), is based on polling by the AP. It is rarely implemented in reality due to unresolved problems such as uncontrolled transmission time of polled MSs. The distributed and contention-based control leads to weak QoS support capability. To achieve better QoS, 802.11e develops new approaches by means of MAC enhancements. The above two access modes are improved with service differentiation. Also, admission control and bandwidth reservation (scheduling) are considered to support multimedia applications with stringent QoS requirements. Various QoS mechanisms for 802.11 are explored in [9]. Nevertheless, WLANs still cannot be expected to support the same level of QoS as 3G networks. Considering the differences of 3G networks and WLANs in QoS provisioning and the aforementioned overlaying structure, different services can be admitted to either the cellular network or the WLAN according to their traffic characteristics and QoS requirements. For example, real-time services such as voice telephony can be well carried by the cellular network to satisfy the strict delay requirement, whereas the delay-tolerant data traffic can be admitted to the WLAN to enjoy the high throughput. It is also an important issue to maintain consistent or smoothly adapted QoS during vertical handoff (i.e., handoff between the cellular network and the WLAN).

D. Interworking Scenarios

Currently, WLANs may be owned by a 3G operator, a commercial WLAN operator, an authority of a public hotspot (e.g., airports or property management corporations), or a business enterprise for internal use [10]. The interworking mechanisms are directly related to the ownership of WLANs. More exposure is possible to integrate the WLANs owned by the cellular operator itself. The objective penetration level between the 3G networks and WLANs and the provisioned services lead to different requirements for the interworking mechanisms. The services supported in future integrated 3G/WLAN networks are envisioned in [11]. However, the integration of the two technologies will be a gradual development process. From the perspective of a 3G operator, a step-wise approach is proposed in [12], which defines six interworking scenarios with each scenario specifying an incremental set of service and operational features. The first scenario only requires common billing and customer care. In the second scenario, a 3G subscriber roaming to the WLAN is authenticated and charged by its 3G home network. Only IP access service via the WLAN is provided to the roaming user. In the third scenario, the 3G packet-switched services are also open to users attached to the WLAN, such as MMS, WAP service, IP multimedia and location-based services. The fourth and fifth scenarios improve the third scenario with higher requirement for service continuity. The sixth scenario allows the access to the 3G circuit-switched services (such as conventional voice calls) via WLANs, and supports seamless mobility across 3G and WLANs. So far, a lot of research in the literature focuses on the four scenarios in the middle.

III. 3G/WLAN INTERWORKING ARCHITECTURES AND MECHANISMS

The standardization for 3G/WLAN interworking is well in progress by 3GPP and 3GPP2. The high-level interworking requirements, architecture and procedures (e.g., network selection, authentication, charging, etc.) have been specified in [13]. In [14] the integration architecture is classified into two categories according to the inter-dependence between the two access networks, i.e., tight coupling and loose coupling. Different interworking mechanisms are developed to support mobility, security, QoS, charging and billing in the tightly-coupled and loosely-coupled heterogeneous networks.

A. Tightly Coupled and Loosely Coupled

Integration

In the tight-coupling architecture, the WLAN is connected to the 3G core network as one 3G radio access network. Figure 1 illustrates a simplified architecture for interworking UMTS/GPRS networks with 802.11 WLANs. The lines with tags \a" and \b" are examples for tight coupling with different integration points. Typical tight-coupling architectures include the first architecture proposed in [15] and [16]. We can see that, in this type of integration architecture, the cellular radio is simply replaced with WLAN radio providing equivalent functions. As a consequence, the 3G protocols and existing network infrastructures can be reused. For example, the user roaming across the two domains is based on the mobility management protocols of 3G networks, thus enhancing the inter-domain mobility management capability. However, the high-layer cellular protocols tailored for highly mobile users in hostile outdoor environments may not operate properly for WLANs [17]. The main disadvantages of the tight coupling approach include:

- (i) an interface in 3G core networks exposed to WLAN is required, which is a challenge as the two domains are likely developed and deployed independently by different operators;
- (ii) a large volume of WLAN traffic will go through the 3G core network, possibly making the latter a network bottleneck; and
- (iii) the WLAN needs to have a protocol stack compatible with that of the 3G networks.

On the contrary, for the loose-coupling approach (shown with the line \c" in Figure 1), the WLAN is connected to the cellular core network indirectly through an external IP network such as the Internet. The second architectures proposed in [15] and [16], respectively, and the operator WLAN system in [18] belong to this category. This type of architecture imposes minimal requirements to modify current WLAN standards, and allows for the flexibility and independence of implementing individually different mechanisms within each network. However, the 3G networks may need to be augmented with extra functionalities such as Mobile IP for mobility management and authentication, authorization and accounting (AAA) support. Moreover, as the two domains are separated, the mobility signalling may traverse a relatively long path, thus inducing relatively high handoff latency. Nonetheless, there are various enhancement mechanisms for Mobile IP to reduce the handoff latency, such as regional registration and dynamic home agent (HA) assignment. Overall, the loose coupling is the preferred solution for both the 3G and WLAN communities as it allows the gradual deployment of hotspots with no or little modification on the 3G networks [16] [19] [20]. It is well recognized that the future networks will be all IP networks. Therefore, it is a natural choice to glue 3G networks and WLANs with the pervasive IP technology. In fact, we can see from the later sections that many mechanisms proposed for the interworking follow the de facto standards in the Internet community.

B. Authentication and Authorization Support

It is expected that in a 3G/WLAN integrated network the wireless terminal will be dual-mode, which means that the terminal will be equipped with network interfaces to both 3G networks and WLANs [11]. However, only one subscription is needed with a 3G operator or a WLAN service provider, who has roaming agreements to support the interworking [20]. A lot of research in the literature considers the scenario in which a 3G subscriber is provided WLAN access via independent WLAN systems. Take the interworking of UMTS and 802.11 WLAN as an example. As mentioned in Section 2, 802.11 WLANs adopt 802.1X for access control, based on the EAP and AAA framework. EAP sets no restriction on specific authentication methods, while the AAA protocols, such as remote authentication dial-in user service (RADIUS)1 or its enhanced version Diameter, offer customized attributes in the authentication messages. Considering the flexibility and interoperability requirement, AAA server is introduced in UMTS [21]. By means of EAP-AKA2, the 3G home AAA server authenticates the UMTS subscribers roaming to a WLAN through the UMTS AKA procedure[1] [10] [19] [20]. The authentication information and subscription profile can be retrieved from the home subscriber server (HSS) or HLR. Similar arrangements are implemented in the cdma2000/WLAN interworking.

There are three entities defined in 802.11X, i.e., supplicant, authenticator, and authentication server. It is natural that the wireless terminal and the 3G AAA server will perform the functions of the supplicant and the authentication server, respectively. The authenticator is responsible for relaying EAP frames sent by the supplicant via EAP-over-LAN (EAPOL), and repacketing them into appropriate AAA messages for onward transmission to the authentication server, and vice versa. In the interworking mechanism, the functions of the authenticator can be carried by the AP [15], the WLAN access router, or a separate WLAN AAA proxy [10]. A successful authentication procedure is summarized in Figure 2. In the operator WLAN architecture proposed for legacy GSM/GPRS networks in [18], the access controller actually acts as the access router shown in Figure 1. It provides IP routing for WLANs and also relays authentication messages for the terminal and authentication server. As above, the signalling between the access controller and the authentication server is based on RADIUS, and the SIM-based authentication is executed with the aid of HLR. However, the authentication signalling between the mobile terminal and the access controller is the operator WLAN-specific network access authentication and accounting protocol (NAAP), which encapsulates GSM authentication messages inside IP packets. This actually limits the applicability of the proposed interworking mechanism and is not compatible with current WLAN or IP standards.

The above approaches actually perform the authentication at layer 2. When Mobile IP is supported in the 3G and WLAN networks, the authentication can also be carried out at the network (IP) layer [20]. In the mobile IP mode of cdma2000, PDSN is augmented with foreign agent (FA) functionality, while HA is introduced to maintain location information of MSs and forward IP packets to MSs via tunnelling. For UMTS, the technical report 23.923 (not technical specification yet) [3] outlines the plan to implement Mobile IP in UMTS IP core network through three steps. In the first stage, the FA functionality is added to GGSN, which can support roaming between WLANs and UMTS via Mobile IP. Stage II provides IP mobility for inter-GGSN handoff, while in stage III GGSN and SGSN are further combined to one intelligent GPRS support node (IGSN) similar to the PDSN in cdma2000. With Mobile IP being implemented, the MS roaming to a WLAN can be authenticated with the MIP registration procedure. An authentication example for cdma2000/WLAN interworking is illustrated in Figure 3.

After obtaining a new care-of address (CoA) from the Agent advertisements of FA, the MS initiates the registration procedure by sending an MIP registration request message to the FA, which will contact the 3G home AAA RADIUS server to authenticate the roaming user via mobile-foreign challenge extension and mobile-AAA authentication extension [22]. In the mobile IP mode of cdma2000, neither challenge handshake authentication protocol (CHAP) nor password authentication protocol (PAP) is performed as in the simple IP mode, because an additional AAA traversal will result in longer initial setup time and re-establishment time [23]. After successful authentication by the RADIUS server, the MIP registration request of the MS will be sent to the HA via HA request message. After receiving HA response indicating the registration is validated by the HA, the AAA RADIUS server acknowledges the PDSN/FA with RADIUS access accept. The PDSN/FA will subsequently send an MIP registration reply message to the MS terminating the authentication and registration procedure.

In addition, it is possible to perform the authentication at the application layer. In [20], it is proposed to authenticate the mobile node via a secured Web page for login over a hypertext transport protocol secured (HTTPS) connection. This type of solution has restricted applications since it requires the support of specific services such as HTTP. In summary, by introducing the flexible AAA framework, the 3G-specific authentication mechanisms can be reused. At the same time, independent WLAN service providers can implement their own preferred authentication methods, which are generally the popular standards in the Internet community, such as extensible authentication protocol-message digest number 5 (EAP-MD5) and extensible authentication protocol- transport layer security (EAP-TLS). On the other hand, by coupling the authentication and authorization procedure with mobility management, the signalling procedure can be simplified. However, the authentication and authorization procedure should be efficient enough to minimize its impact on handoff latency. As for higherlayer authentication mechanisms, although more freedom is allowed for underlying technologies, the messaging delay and specific application restriction are their main drawbacks.

a. Mobility Management across 3G Networks and WLANs

Mobility management consists of two aspects, i.e., location management and handoff management. Location management continuously tracks the locations of MSs, while handoff management maintains the on-going connections while switching attachment points. The mobility protocols can operate at different layers such as link, network, transport and application layers [24]. Link-layer mobility support is inherent in each wireless/mobile network, e.g., the GMM in UMTS/GPRS networks. In the tightly-coupled inter- working architecture, the WLANs are integrated as any other 3G radio access networks. As a result, the 3G mobility management mechanism can be reused to support mobility across WLANs and 3G networks [15] [16]. The most popular network layer mobility protocol is Mobile IP [2]. By introducing mobility agents and IP tunnelling, the upper layer applications are enabled

transparency to IP address changes due to user movement. However, the original MIP protocol suffers from the triangle routing problem. The packets have to be first routed to the HA before being tunnelled to the mobile node, while a mobile node sends its packets through a router on the foreign network assuming the routing is independent of source address. When this assumption is not valid, topologically correct reverse tunnels can also be established from the CoA to the HA. As a consequence, an extra delay results from the long route when the visited foreign network is far from the home network. This problem can be solved by route optimization, in which a direct route is established between the mobile node and its correspondent node. In addition, the location registration to the HA may lead to a heavy signalling load when there are a large number of mobile nodes and also a long signalling delay if the foreign and home networks are far apart. Therefore, MIP is more suited for macro-mobility with infrequent movement and often between different administrative domains (inter-domain). There are many MIP variants to support micro-mobility, which refers to mobility within one domain (intra-domain). The handoff latency can be reduced by means of localizing signalling via regional/hierarchical registration (tunnelling-based) or host specific routing (routing-based) [25]. For tunnel-based micro-mobility protocols, there are Mobile IP regional registration, hierarchical Mobile IP and intra-domain mobility management protocol (IDMP). Typical routing-based micro-mobility protocols include cellular IP and HAWAII [26].In the loosely-coupled interworking architectures proposed in [15] and [20], the MIP approach is used for the mobility across the cellular network and WLANs. Even though the loose coupling allows for the flexibility of implementing different individual mobility management schemes within each network, it has become a trend to apply MIP in 3G networks to pave the way for future all-IP networks. A cross-layer solution is proposed in [26] to support mobility in all-IP-based wireless networks with heterogeneous access technologies. The architecture can be applied to the 3G/WLAN interworking case. In the scheme, the aforementioned micro-mobility protocols are used for intradomain mobility, while a cross-layer mobility protocol is proposed for inter-domain mobility. With early detection of possible inter-domain handoff using link layer information, it is enabled to carry out authentication, authorization, and MIP registration before the actual handoff. As a result, the inter-domain handoff delay is reduced to be comparable to an intra-domain handoff delay. Moreover, transport layer mobility protocols can be used to prevent transport-layer applications such as transport control protocol (TCP) connections from being interrupted by IP address changes due to user mobility. One solution in this type of mobility support is TCP-Migrate [27]. Another transport layer scheme proposed in [28] supports UMTS/WLAN vertical handoff via stream control transmission protocol (SCTP). Although performing mobility management at the transport layer provides network- independence, it requires more functions carried by end systems and also many modifications to current end systems.

Application layer mobility protocols can provide mobility support independent of underlying wireless access technologies and network layer protocols [28]. Hence, they offer another alternative for mobility management in the 3G/WLAN interworked networks. A typical representative of application layer mobility protocols is session initiation protocol (SIP) [24], which supports pre-call and mid-call mobility with the aid of application-specific redirect server. Although the application layer solution introduces less modification to existing protocol stacks and infrastructures of 3G networks and WLANs, longer handoff latency may be incurred due to the high-layer messaging. In [29], a hybrid multilayer mobility management scheme is proposed. To handle macromobility, it uses Mobile IP for non-real time services and SIP for real-time services.

Existing protocols such as cellular IP, HAWAII and hierarchical Mobile IP are employed to support micro-mobility.

AAA context transfer is further introduced to avoid additional delay induced by the AAA security procedure during handoff.

In summary, for mobility management in 3G/WLAN integrated networks, accurate location information of MSs needs to be maintained, or the MSs can be effectively located by paging, so that the MSs can move freely across the two domains. Moreover,

the vertical handoff is required to be fast, smooth and seamless [30]. Especially for real-time services, the vertical handoff latency needs to satisfy strict delay bounds. Smooth handoff requires minimized packet loss during handoff, while the perceptible service interruption should be minimized to achieve the objective of seamless handoff. In a 3G/WLAN interworking environment involved with multiple domains, when enhanced with micromobility protocols, the Mobile IP approach is a good solution and can well meet the above requirements, especially when hierarchical layered mobility management architecture is applied. The performance of Mobile IP can be further improved with cross-layer design. Architecturally, the network layer is the right place to handle mobility [24]. The higher-layer approaches such as the transport-layer and application-layer mobility management schemes have the advantage of retaining the network interfaces unchanged. This feature applies well to the heterogeneous nature of 3G/WLAN integrated networks. Nonetheless, compared with the network-layer approaches, these higher-layer solutions are relatively immature and incomplete.

D. QoS Provisioning Architecture

The next-generation wireless networks are required to support multimedia services, which is actually a driving force for the 3G/WLAN interworking. Multimedia applications usually feature more bandwidth-demanding and stringent QoS requirements. We have mentioned in Section 2.3 the differences between 3G networks and WLANs in QoS support at the physical layer and MAC layer. This heterogeneous nature results in technical challenges in end-to-end QoS provisioning for a variety of services in an integrated 3G/WLAN network.

Figure 4 demonstrates the various scenarios requesting end-toend services. It can be seen that different networks may be involved along the end-to-end path, e.g., 3G radio access networks, 3G core networks, WLAN access networks, IP backbone, and the wired Internet. In such a case, the following aspects should be well addressed in the QoS provisioning architecture [31]:

- QoS attributes The traffic characteristics and end-to-end QoS requirements need to be clearly specified by the QoS attributes. These QoS attributes should be independent of the underlying access technologies, easily interpreted and mapped to service requirements for the networks involved along the end-to-end path;
- QoS signalling protocols The QoS signalling protocols are responsible for QoS negotiation, resource reservation, etc. Although they are not required to be identical across different network domains, it is essential for the QoS attributes to be distributed via signalling along the end-to-end path without modification and being kept consistent;
- QoS mechanisms To satisfy the QoS requirements presented to each network segment along the end-to-end path, proper local QoS mechanisms such as traffic control and scheduling are evoked to provide the requested QoS. Distinct QoS mechanisms are employed in 3G networks and WLANs, which results in different levels of QoS support capability.

A good overview of QoS provisioning in 3G networks is given in [32]. For UMTS, a layered QoS architecture is specified [33]. The end-to-end bear required to meet the QoS objectives is decomposed into a concatenation of bear services, which are provided by the different network segments involved. There are four traffic classes supported in UMTS, i.e., conversational, streaming, interactive and

background, with QoS attributes defined for each class, such as maximum bit rate, transfer delay, and traffic handling priority. As for the QoS signalling, the UMTS packet data protocol (PDP) context signalling mechanism is used for the bear services within UMTS. Diverse QoS schemes can be employed in the air interface, radio access networks and core networks to achieve the QoS objectives, e.g., power control, bandwidth scheduling, and transport technologies.

As mentioned in Section 3, IP technologies offer a common platform to integrate WLANs with 3G networks. In UMTS, the gateway node to external packet networks (i.e., GGSN) is required to provision an IP bearer service manager for interworking with

external IP networks. For IP networks, there are two major OoS architectures, i.e., integrated services (IntServ) and differentiated services (DiffServ). In the IntServ architecture, with the aid of resource reservation protocol (RSVP), resources are reserved along the end-to-end path to satisfy the QoS requirements for each flow. However, the per-flow soft states must be maintained end to end and refreshed periodically to hold the reservation. As a result, this model is not well scalable for large networks with a large population. On the other hand, DiffServ does not need a separate signalling protocol and it only exerts provided to different classes. In addition, the edge nodes can be augmented with functionalities such as traffic policing and admission control to prevent unacceptable QoS degradation due to saturation. As such, DiffServ only provides relative QoS to aggregate traffic flows instead of absolute QoS guarantee to individual flows as IntServ. Hence, although DiffServ has good scalability, only coarse QoS provisioning is enabled treatment according to per-hop behavior (PHB) defined for each service class. PHBs are actually implemented by means of buffer management and packet scheduling mechanisms -. In this way, service differentiation can

Because both IntServ and DiffServ models are proposed for wired IP networks, unique characteristics of wireless links and user mobility should be taken into account for extending them to wireless/mobile networks. The extension mechanisms of RSVP and DiffServ for IP-based wireless mobile networks are surveyed in and -, respectively. Especially, in the case of integrated 3G and WLANs networks, when vertical handoff occurs between the 3G networks and WLANs, the bear services offered by the two access technologies provide different QoS support. These changes should be considered in the OoS context transfer and service reestablishment after handoff control for aggregate traffic. At the network edge, the traffic flows are classified into one of several classes according to the service level specification (SLS) between the user and the access networks and between different network domains [34]. The type of service (TOS) field of IP packets is marked with the differentiated service code point (DSCP) indicating the service class. Then the packets will receive prioritybased

In an IP QoS model combining IntServ and DiffServ is proposed for heterogeneous wireless networks. In the wireless access networks, IntServ architecture is applied, since the traffic volume is relatively low and scalability will not be a severe problem. Also, the scarce wireless bandwidth can be efficiently utilized by properly allocating resources to each traffic flow. On the other hand, the core networks adopt the DiffServ architecture for scalability purpose. At the edge of the sender's access networks, the RSVP signalling for end-to-end resource reservation is mapped to DiffServ PHBs, and the RSVP messages are tunnelled through the DiffServ domains till the receiver end. The DiffServ approach for core networks is further justified by the bandwidth overprovisioning in current core networks and the relatively low cost of upgrading transmission capability . A similar idea is used in the QoS architecture proposed in [34], where RSVP is slightly modified. Instead of exchanging the RSVP PATH and RSVP RESV messages end to end from the sender to the receiver, the RSVP signalling is localized within the access networks via proxies. This facilitates the reservation path repair after handoff. In addition, it is proposed in that multiprotocol label switching (MPLS) can be introduced in the DiffServ domain. Due to the intrinsic QoS provisioning capability, MPLS can facilitate scalability by means of flow aggregation and at the same time guarantee individual QoS without the need to maintain per-flow awareness along the path. To achieve better QoS support, mobility management and QoS provisioning can be coupled in the architecture design. In, mobility and QoS architecture is proposed for all-IP based 4G wireless networks. It takes advantage of the two-layer hierarchical structure of IDMP mobility management to reduce the resource reservation delay due to RSVP reestablishment after handoff. The per-flow RSVP reservation is terminated at a fixed anchor subnet agent, while a forwarding chain is used to track the host mobility within a domain. As such, the reservation messages can be effectively restricted within a local scope. Another IP-based QoS architecture is proposed in for 4G networks

supporting multiple accesses and multiple service providers. In this architecture, the end-to-end QoS support is integrated with mobility and authentication, authorization, auditing and charging (AAAC). In each network domain, an AAAC system is used to handle the network access control, while at least one QoS broker (QoSB) manages the access network resources. The common open policy control (COPS) protocol is used for information exchanging between QoSBs, edge routers and AAAC systems, because in this architecture QoSBs and AAAC systems act as policy decision points (PDPs) while edge routers act as policy execution points (PEPs).

In, a policy-based multidomain QoS architecture is proposed to provide consistent QoS control over an integrated UMTS and WLAN system. The architecture varies with different interworking scenarios such as interworking with a UMTS operator's WLAN, a WLAN shared by multiple operators, or a WLAN of a UMTS operator's customer. This policy-based type of approach views the development of proper QoS architecture for 3G/WLAN interworking from a new perspective. However, there are still some problems that need more careful considerations, such as fastening policy negotiation, and securing information exchanging between interconnected policy entities.

IV. CONCLUSION

In this article, we have introduced the complementary characteristics of 3G networks and WLANs and the motivations for their interworking. The differences of the two wireless technologies lead to many challenging issues that need to be addressed in the interworking. To achieve seamless roaming, enhanced security, and consistent QoS, there are a variety of interworking mechanisms proposed from different perspectives. In particular, we examine the typical solutions for user authentication, mobility management, and QoS provisioning in the 3G/WLAN interworked networks. It is observed that IP technologies play an important role in gluing the two heterogeneous technologies. The demand for ubiquitous Internet access further strengthens this convergence trend. For instance, the well-known AAA framework, Mobile IP protocol, IntServ/RSVP and DiffServ QoS architectures are considered as promising approaches for interworking 3G networks and WLANs. However, most current research on 3G/WLAN interworking focuses on relatively high-level architecture issues. Although the local QoS mechanisms implemented in the 3G networks and WLANs can provide certain levels of QoS support, proper resource allocation schemes are needed to utilize the integrated resources effectively and efficiently. In , we investigate the resource allocation problem in 3G/WLAN integrated networks. In the proposed resource allocation scheme, we try to take advantage of the characteristics of this heterogeneous network with multiple services, so that better QoS assurance and higher utilization can be achieved. In, a radio resource scheduling scheme is proposed for a 3G/WLAN coupled

network. Based on the traffic characteristics, QoS requirements, and network states, a joint scheduler prioritizes and splits the traffic over the two networks. A certain amount of resources in each network is allocated to the corresponding traffic sub stream to achieve the target QoS. A higher utilization is available due to a larger trunking gain and effective manipulation of the overall resources of the two networks. However, these are only some initial research attempts for the resource allocation problem. It remains to be an open issue that needs more in-depth investigation.

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