



## Cryptography of a Binary Image using a Modified Hill Cipher

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**Abstract:** In this paper, we have used a modified Hill cipher for encrypting a binary image. Here, we have illustrated the process of encryption by considering a couple of examples. The security of the image is totally achieved, as the encrypted version of the image, does not reveal any feature of the original image.

**Keywords:** Cryptography, Cipher, Binary image, Encrypted image, Modular arithmetic inverse.

### I. INTRODUCTION

In a recent investigation [1, 2], we have modified the Hill cipher by developing an iterative procedure. In [1], we have multiplied the plain text matrix  $P$  with a key matrix  $K$  on both the sides of  $P$ , while in [2] we have used  $K$  on one side and  $K^{-1}$  on the other side of  $P$  as multiplicands. The process is strengthened by using a function called Mix ( $P$ ), at any stage of the iterative process. It is further supplemented with the XOR operation between the plain text  $P$  and the key  $K$ . In this, the key is containing 32 decimal numbers, which are in the interval  $[0 - 255]$ , and the modular arithmetic inverse of the key, represented in the form of a  $16 \times 16$  matrix, is obtained by using mod 2.

In the present paper, our objective is to develop a block cipher, and to use it for the cryptography of a binary image. Here, we have taken a key containing 32 decimal numbers [1, 2], and generated a key matrix of size  $32 \times 32$  in a special manner (discussed later), and applied it in the cryptography of a pair of binary images.

In Section 2, we have developed a procedure for the cryptography of a binary image. In Section 3, we have used a pair of examples and illustrate the process. Finally, in Section 4, we have drawn conclusions from the analysis.

### II. DEVELOPMENT OF A PROCEDURE FOR THE CRYPTOGRAPHY OF A BINARY IMAGE

Consider a binary image whose gray level values can be represented in the form of a matrix given by

$$P = [P_{ij}], \quad i = 1 \text{ to } n, j = 1 \text{ to } n. \quad (2.1)$$

Here, each  $P_{ij}$  is 0 or 1, where 1 corresponds to black and 0 corresponds to white.

Let us choose a key  $k$ . Let it be represented in the form of a matrix given by

$$K = [K_{ij}], \quad i = 1 \text{ to } n, j = 1 \text{ to } n, \quad (2.2)$$

where each  $K_{ij}$  is either 0 or 1.

$$\text{Let } C = [C_{ij}], \quad i = 1 \text{ to } n, j = 1 \text{ to } n \quad (2.3)$$

be a matrix, obtained on encryption.

The process of encryption and the process of decryption, which are quite suitable, for the problem on hand, are given in Fig. 1.

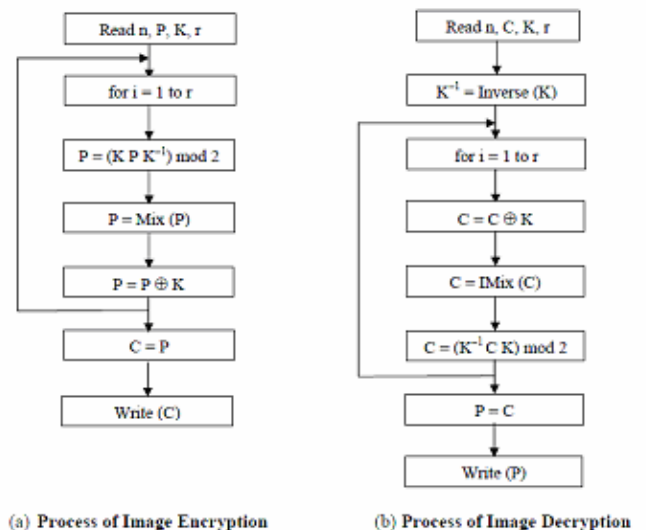


Figure 1: Schematic diagram for the cryptography of a image

Mix ( ) is a function used for mixing thoroughly the binary bits arising in the process of encryption at each stage of iteration. IMix ( ) is a function which represents the reverse process of Mix ( ). For a detailed discussion of these functions, and the algorithms involved in the processes of encryption and decryption, we refer to [1].

### III. ILLUSTRATION OF THE CRYPTOGRAPHY OF AN IMAGE

Let us take a key  $k$  consisting of 32 decimal numbers. This is given by

$$k = \left\{ \begin{array}{cccccccccccccccc} 131 & 31 & 18 & 59 & 254 & 126 & 113 & 97 & 127 & 167 & 76 & 116 & 111 & 159 & 245 & 159 \\ 175 & 50 & 236 & 107 & 235 & 74 & 47 & 20 & 190 & 80 & 242 & 139 & 175 & 164 & 187 & 158 \end{array} \right\} \quad (3.1)$$

Let us write  $k$  in the form of a matrix by placing the first two numbers 131 and 31, in their binary form, in the first row, the next two numbers 18 and 59, in the second row, and so on.

Thus we get a matrix, denoted as  $Q$ , in the form

$$Q = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 1 & 1 & 1 & 0 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 0 \\ 0 & 1 & 1 & 1 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 1 \\ 0 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 0 & 1 & 0 & 1 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 0 \\ 1 & 1 & 1 & 0 & 1 & 0 & 1 & 0 & 0 & 1 & 1 & 0 & 1 & 0 & 1 & 1 \\ 1 & 1 & 1 & 0 & 1 & 0 & 0 & 1 & 0 & 1 & 0 & 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 1 & 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 1 & 0 & 0 & 0 & 1 & 0 & 1 & 1 \\ 1 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 0 & 1 & 0 & 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 1 & 1 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 1 & 1 & 1 & 0 \end{bmatrix} \quad (3.2)$$

Now, in order to have a larger key, for convenience, let us take a key  $K$  in the form

$$K = \begin{bmatrix} Q & R \\ S & U \end{bmatrix} \quad (3.3)$$

where  $U = Q^T$ , in which  $T$  denotes the transpose, and matrices  $R$  and  $S$  are obtained from  $Q$  and  $U$  respectively, by adopting the following procedure.

To obtain R, we interchange the first row of Q with its 16<sup>th</sup> row. Similarly, we interchange the second row of Q with its 15<sup>th</sup> row. We continue this process till we exhaust all the remaining rows of Q. Thus we get R. In the same manner, we obtain S from U. Now a circular rotation through 4 rows is applied in the downward direction (i.e., the first row becomes the fifth row and the rest of the rows move in the same direction past 4 rows) on U, and hence we get the new U. Thus, we have

[illegible]

The afore mentioned operations are performed for

- (1) enhancing the size of the original key matrix ( $16 \times 16$ ) to  $32 \times 32$ , and

- (2) obtaining the modular arithmetic inverse of  $K$ , in a trial and error manner.

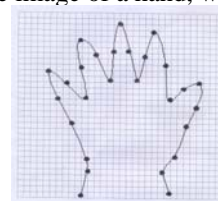
The modular arithmetic inverse is obtained as

[illegible]

From (3.4) and (3.5), we can readily find that

$$\mathbf{K} \mathbf{K}^{-1} \bmod 2 = \mathbf{K}^{-1} \mathbf{K} \bmod 2 = \mathbf{I}. \quad (3.6)$$

Let us consider the image of a hand, which is given below.



### Fig. 2. Image of a Hand

This image can be represented in the form of a binary matrix  $P$  given by

[illegible]

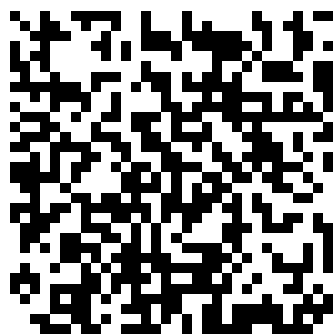
where 1 denotes black and 0 denotes white.

On adopting the iterative procedure given in Fig. 1, we get the encrypted image C in the form

[illegible]

On using (3.4), (3.5), (3.8), and the procedure for decryption (See Fig. 1.(b)), we get back the original binary image P, given by (3.7).

From the matrix C, on connecting each 1 with its neighbouring 1, we get an image which is in a zigzag manner (See Fig. 3).



**Fig. 3. Encrypted image of the hand**

It is interesting to note that, the original image and the encrypted image differ totally, and the former one, exhibits all the features very clearly, while the later one does not reveal anything.

Now let us consider another image which is consisting of the upper half of a person. This is shown Fig. 4.



### Fig. 4. Image of a Person

This image can be represented in the form of a matrix P, containing binary bits, as shown below.

[illegible]

On using (3.4), (3.5), (3.9), and encryption algorithm given in Fig. 1.(a), we get the encrypted image C, given by

[illegible]

On applying the process of decryption given in Fig. 1.(b), we get the original binary image given in (3.9).

On comparing the matrices of P and C, we find that, the 1s in P are scattered when P is transformed into C. On connecting 1s in C, we get an image as shown in Fig. 5.



**Fig. 5. Encrypted image of a Person**

Here, we also notice that, the features of the original image are totally lost, when the original image is transformed into the encrypted one.

#### IV. CONCLUSIONS

In this analysis, we have made use of a modified Hill cipher for encrypting binary images. In this, we have illustrated the procedure by considering a pair of examples: (1) the image of a hand, and (2) the image of the upper half of a person.

Here, we have noticed that, the encrypted images are totally different from the original images, and the security of the images is completely enhanced, as no feature of the original images can be traced out in any way from the encrypted images.

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This analysis can be extended for the images of signatures and thumb impressions.

#### V. REFERENCES

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